Locality and the Fast File System

Chapter 41

Previously in CS212…

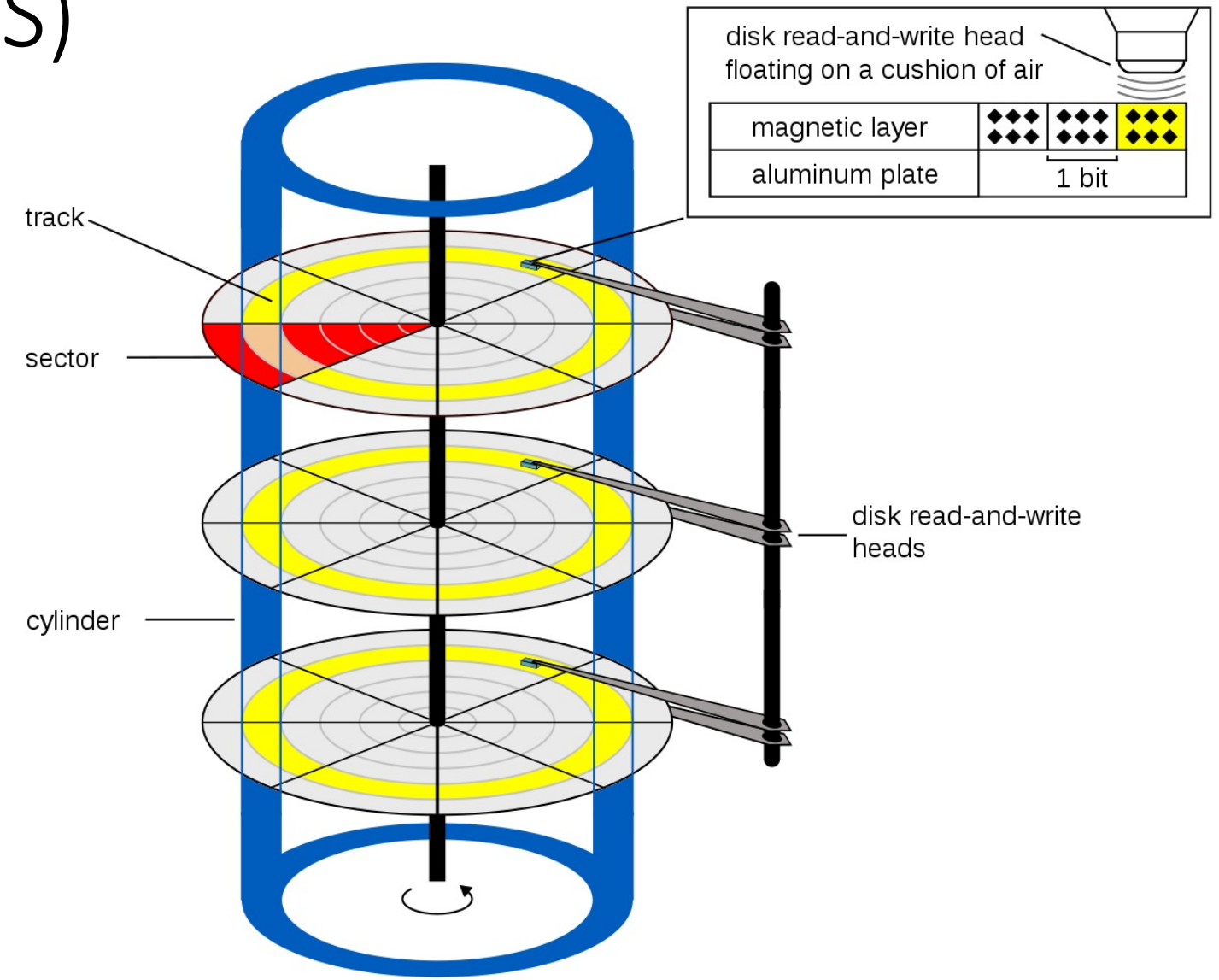
- We looked at the basics of a file system implementation
- Disks are broken down into blocks
- We use part of the disk for holding data, and other portions for metadata
  - Inodes for block free list and file/directory location metadata
  - SUPERBLOCK for metadata about the filesystem
- However, we have a problem….

Poor Performance

- The vsfs example and the old UNIX filesystem don’t store things with locality in mind
- Data is very far from the inodes which causes expensive seek operations
- Small block sizes minimized internal fragmentation
  - External fragmentation can cause files to be spread across multiple non-consecutive blocks
  - Mechanical HDDs benefit from defragmentation tools

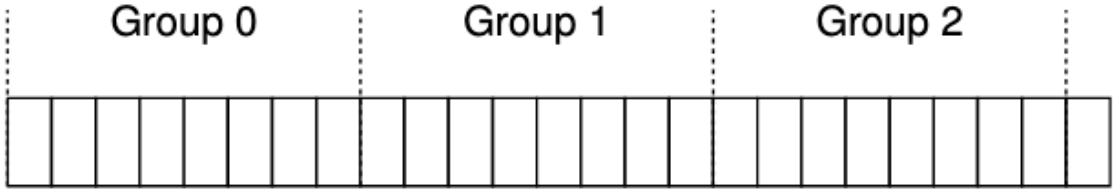
Fast File System (FFS)

- Built on top of the standard interface
  - open(), read(), write(), close()
- A **cylinder** is the set of tracks in the same position across all patter surfaces
- FFS collects consecutive cylinders into a **cylinder group**



Logical Organization

- The HDD doesn't share information about the geometry of the HDD
  - Only block addresses
- Modern file systems logically organize the drive into **block groups**
  - Each group is a consecutive portion of the disk address space



- The important thing is that data stored in the same group will not result in long seeking across the disk
- Each group keep track of its own file system structures

Allocating Files and Directories

- Keep related stuff together
  - Allocate data blocks for a file in the same group as its inode
  - Place all files that are in the same directory in the cylinder group of the directory they are in
- Assume we store /a/c, /a/d, /a/e, /b/f

Cylinders with directory locality

group	inodes	data
0	/	/
1	acde	accddee
2	bf	bff
3		
4		
5		
6		
7		

Vs.

Cylinders without directory locality

group	inodes	data
0	/	/
1	a	a
2	b	b
3	c	cc
4	d	dd
5	e	ee
6	f	ff
7		

Large-File Exception

- A large file could take up all the space in a group
  - This means we might not be able to store other files in the same directory within the same group

```
group inodes     data
  0   /a-------- /aaaaaaaaa aaaaaaaaaa aaaaaaaaaa a---------
  1   ---------- ---------- ---------- ---------- ----------
  2   ---------- ---------- ---------- ---------- ----------
```

- We instead use up the direct blocks first and use the indirect pointers to store the remaining content in other block groups
  - Likely with less utilization

```
group inodes     data
  0   /a-------- /aaaaa---- ---------- ---------- ----------
  1   ---------- aaaaa----- ---------- ---------- ----------
  2   ---------- aaaaa----- ---------- ---------- ----------
  3   ---------- aaaaa----- ---------- ---------- ----------
  4   ---------- aaaaa----- ---------- ---------- ----------
  5   ---------- aaaaa----- ---------- ---------- ----------
  6   ---------- ---------- ---------- ---------- ----------
```

- Does have performance issues, but we can limit it with larger chunks
  - A chuck is just a unit of how much data we read/write from the disk

Other FFS Features

- Sub-blocks to hold small files
  - Reduce internal fragmentation
  - Mostly avoided by having the library buffer the data and write out when it's large enough
- Parameterization
  - Skip blocks when writing to consecutive data to about "missing" data on a rotation
  - Modern drives have a track buffer to hold a track and reach from this cache on subsequent reads for that track

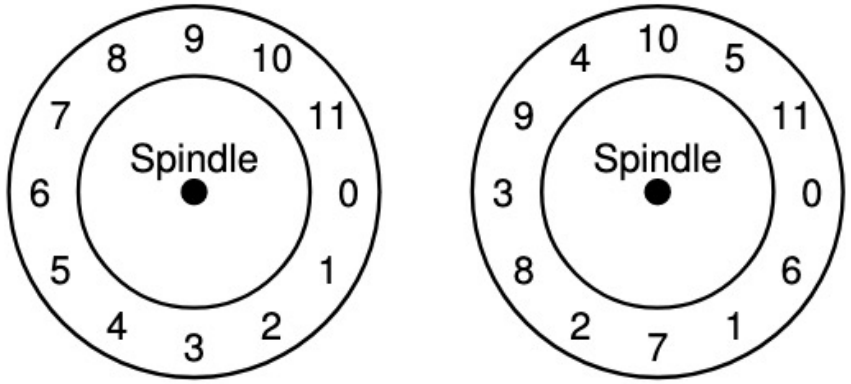


Figure 41.3: **FFS: Standard Versus Parameterized Placement**

Next Time

- We talk about how we handle situations when things go wrong…