

Potential Topic Areas

Games and Graphics

Procedural terrain generation

Optimizing something

Particle systems

2D to 3D model generation

Text Processing

File parsing

Computer Architecture

RISC-V

More Theoretical Topics

Automated Theorem Proving

SAT solvers

Turing machines

Cybersecurity

Cryptography

Robotics

AI / Machine Learning / Data Science

Playing a game

Algorithmic trading

Web Development

Networking

Software Engineering

Front end

Backend

Cloud computing

Divide and Conquer

- Recursive approach
- Break a problem into subproblems that are smaller instances of the same problem to create a solution to the original problem
- Steps
 - Divide into subproblems
 - Conquer the subproblems by solving them recursively, or if the subproblem sizes are small enough, solve them in a straightforward manner (base case of recursion)
 - Combine the solutions to subproblems into the solution to the original problem

MERGE-SORT(A, p, r)

if $p < r$

$q = \lfloor (p + r) / 2 \rfloor$

 MERGE-SORT(A, p, q)

 MERGE-SORT($A, q + 1, r$)

 MERGE(A, p, q, r)

 // check for base case

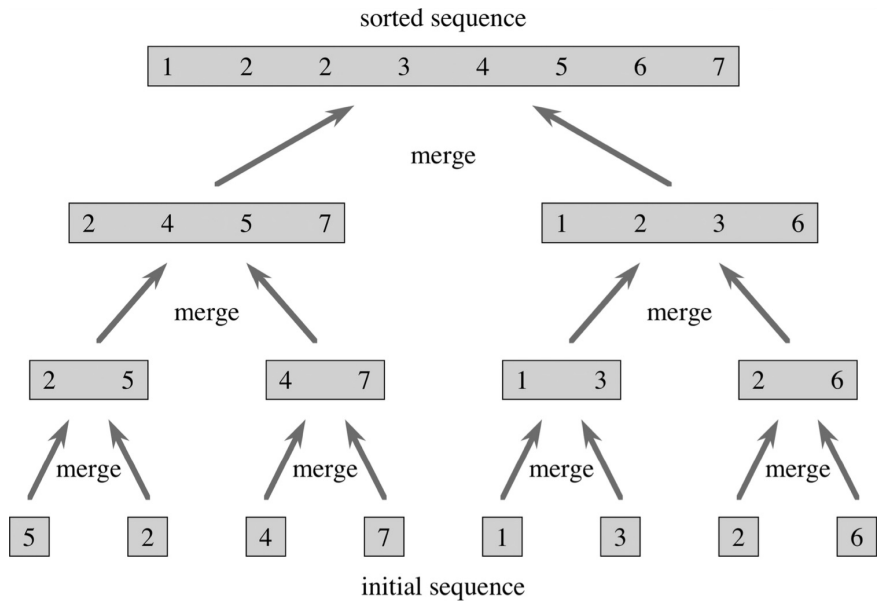
 // divide

 // conquer

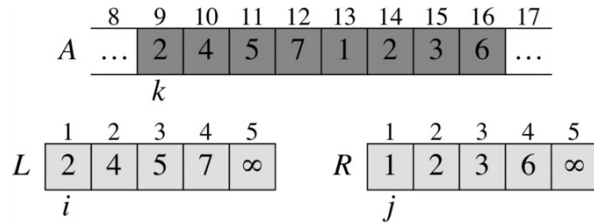
 // conquer

 // combine

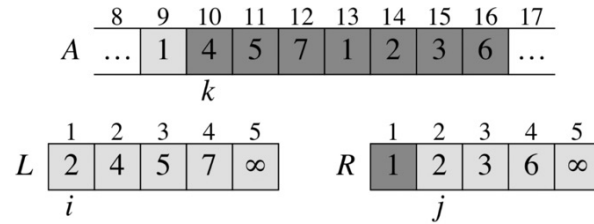
Sorting subarray $A[p..r]$



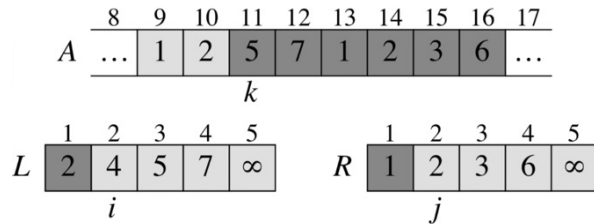
Merge sort requires using extra storage during the merge step



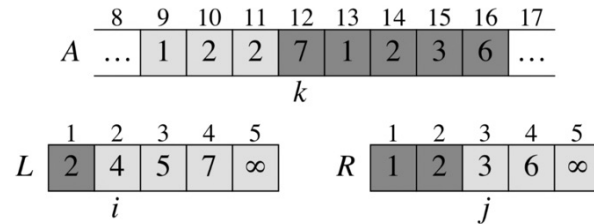
(a)



(b)



(c)



(d)

	8	9	10	11	12	13	14	15	16	17
A	...	1	2	2	3	1	2	3	6	...

k

	1	2	3	4	5
L	2	4	5	7	∞

i

	1	2	3	4	5
R	1	2	3	6	∞

j

(e)

	8	9	10	11	12	13	14	15	16	17
A	...	1	2	2	3	4	2	3	6	...

k

	1	2	3	4	5
L	2	4	5	7	∞

i

	1	2	3	4	5
R	1	2	3	6	∞

j

(f)

	8	9	10	11	12	13	14	15	16	17
A	...	1	2	2	3	4	5	3	6	...

k

	1	2	3	4	5
L	2	4	5	7	∞

i

	1	2	3	4	5
R	1	2	3	6	∞

j

(g)

	8	9	10	11	12	13	14	15	16	17
A	...	1	2	2	3	4	5	6	6	...

k

	1	2	3	4	5
L	2	4	5	7	∞

i

	1	2	3	4	5
R	1	2	3	6	∞

j

(h)

	8	9	10	11	12	13	14	15	16	17
A	...	1	2	2	3	4	5	6	7	...

k

	1	2	3	4	5
L	2	4	5	7	∞

i

	1	2	3	4	5
R	1	2	3	6	∞

j

(i)

MERGE(A, p, q, r)

Merges the sorted subarrays $A[p..q]$ and $A[q+1..r]$

$$n_1 = q - p + 1$$

$$n_2 = r - q$$

let $L[1..n_1 + 1]$ and $R[1..n_2 + 1]$ be new arrays $\Theta(n)$

for $i = 1$ to n_1

$$L[i] = A[p + i - 1] \quad \Theta(n)$$

for $j = 1$ to n_2

$$R[j] = A[q + j] \quad \Theta(n)$$

$$L[n_1 + 1] = \infty$$

$$R[n_2 + 1] = \infty$$

$$i = 1$$

$$j = 1$$

for $k = p$ to r

if $L[i] \leq R[j]$ $\Theta(n)$

$$A[k] = L[i]$$

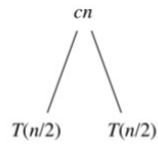
$$i = i + 1$$

else $A[k] = R[j]$

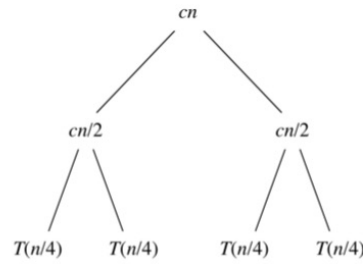
$$j = j + 1$$

Merging is $\Theta(n)$

$T(n)$

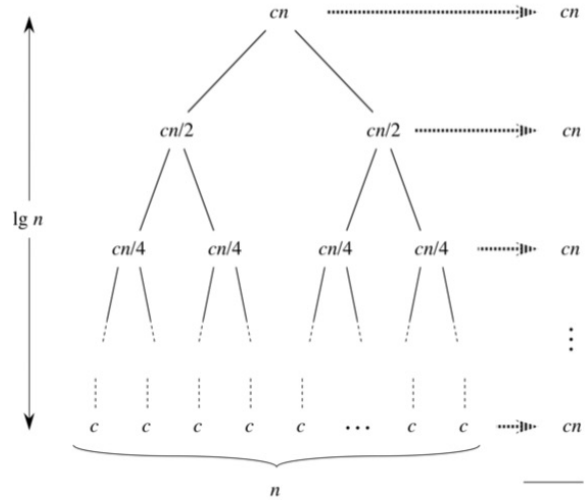


(a)



(b)

(c)



(d)

Total: $cn \lg n + cn$